

Efficient Compilation of Circuit Equations of Switched Networks

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Abstract

The algorithm of finding the set of independent circuit variables of the networks with periodically operated switches is proposed. Our method is based on so-called voltage and current coefficients. As a result, number of circuit equations is considerably reduced in comparison to the classical modified nodal approach.

1 BASIC CONCEPTION

For computer analysis of switched network, the number of circuit equations plays important role. Large number of equations leads to the time- and memory- expensive simulation. Classical modified nodal approach compiles additional equations for switches and transform cells and number of equations grows unsatisfactorily.

The two-graph modified nodal approach (MNA) [1,2] has been used as the basic method. To create minimum number of equations, our procedure is as follows:

- We find the set of independent nodal variables in each switching phase:

$$\begin{aligned} \mathbf{v}_1 &= [v_1^1, v_1^2, \dots, v_1^{n_1}]^T \dots \text{phase 1,} \\ \mathbf{v}_2 &= [v_2^1, v_2^2, \dots, v_2^{n_2}]^T \dots \text{phase 2,} \end{aligned} \quad (1)$$

where n_1 and n_2 are numbers of independent variables in phase 1 and phase 2, respectively. The algorithm how to obtain that sets will be described in following chapter.

- It is possible to create the set of differential equations of MNA in the operator domain for independent variables:

$$\begin{aligned} \underbrace{(\mathbf{G}_1^M + s\mathbf{C}_1^M)}_{\mathbf{Y}_1^M} \mathbf{V}_1 &= \mathbf{D}_1 \mathbf{W}_1 + \mathbf{C}_{12}^M \mathbf{v}_2 (kT^-), \\ \underbrace{(\mathbf{G}_2^M + s\mathbf{C}_2^M)}_{\mathbf{Y}_2^M} \mathbf{V}_2 &= \mathbf{D}_2 \mathbf{W}_2 + \mathbf{C}_{21}^M \mathbf{v}_1 (kT + T_1^-), \end{aligned} \quad (2)$$

where $\mathbf{G}_1^M$ and $\mathbf{G}_2^M$ are modified conductance matrices of dimensions $(n_1 \times n_1)$ and $(n_2 \times n_2)$,

$\mathbf{C}_1^M$ and $\mathbf{C}_2^M$ are modified reactance matrices of dimensions $(n_1 \times n_1)$ and $(n_2 \times n_2)$,

$\mathbf{C}_{12}^M$ and $\mathbf{C}_{21}^M$ are modified transitive matrices of dimensions $(n_1 \times n_2)$ and $(n_2 \times n_1)$,

$\mathbf{D}_1$ and $\mathbf{D}_2$ are certain incidence vectors of dimensions $(n_1 \times 1)$ and $(n_2 \times 1)$,

$\mathbf{W}_1$ and $\mathbf{W}_2$ are Laplace transformations of input signals in phases 1 and 2,

T_1 is length of phase 1.

Equations (2) have key significance for computer analysis. It is possible to compile them for arbitrary linear switched network no matter whether the consistent initial conditions occur in the time instances between two switching phases [3]. In practice, the simulator will produce correct results also in case of discontinuously changed nodal variables.

In following part, we describe the algorithm how to find independent circuitvariables.

2 FINDING OF THE INDEPENDENT CIRCUIT VARIABLES

One way how to find independent circuit variables is given in [4] as a method of current and voltage graphs. Number of compiled equations will be equal to the number of independent circuit variables. To find the set of independent variables, we have to perform transformation and reduction of variables even before the compilation of circuit equations.

phase	node	consecutive numbering * **		transformation due to				arrangement * **	
				switches * **		buffers * **			
1	1	1	1	1	1	1	1	1	1
	2	2	2	1	1	1	1	1	1
	3	3	3	3	3	3	3	2	2
	4	4	4	3	3	3	3	2	2
	5	5	5	3	3	3	3	2	2
	6	6	6	6	6	6	0	3	0
	7	7	7	0	0	0	0	0	0
	8	8	8	8	8	6	8	3	3
2	1	1	1	0	0	0	0	0	0
	2	2	2	0	0	0	0	0	0
	3	3	3	0	0	0	0	0	0
	4	4	4	4	4	4	4	1	1
	5	5	5	5	5	5	5	2	2
	6	6	6	6	6	5	0	2	0
	7	7	7	6	6	5	0	2	0
	8	8	8	5	5	5	5	2	2

Tab.1. Finding of the set of independent variables of the filter in Fig.1.

4 ALGORITHM

It is evident that the algorithm of finding the independent circuit variables will have following parts:

- transformation due to switches,
- skipping the isolated nodes,
- transformation due to nonregular elements,
- rising numbering of the coefficients.

First of all, two matrices U_1 and U_2 [0..N,0..1] are declared, where N is the total number of nodes except for the reference node. The voltage/current coefficients of the k -th node will be recorded to the first/second column and k -th row of matrix U_1 for the phase 1 and U_2 for the phase 2. The result of the algorithm will be as follows:

voltage coefficient of the k -th node: $k^* = U_1(k,0)$ for the phase 1, $U_2(k,0)$ for the phase 2;

current coefficient of the k -th node: $k^{**} = U_1(k,1)$ for the phase 1, $U_2(k,1)$ for the phase 2.

Before starting of the algorithm, both matrices are filled so that both coefficients in the k -th row are k (see Tab. 1, the column "consecutive numbering"). It corresponds to the initial state without transformations.

(a) Transformation Due To Switches

In this stage of the algorithm, the voltage and current coefficients of node are equal. If the connection of nodes a and b is indicated in given phase, it is necessary to perform following operations:

- To read values of coefficients $S1 = a^* = a^{**}$ and $S2 = b^* = b^{**}$ from the corresponding matrix U . Comparing both coefficients and their prospective interchange so that $S1 \leq S2$. Then to replace higher coefficient by the lower on the given position.
- To rewrite all coefficients $S2 \rightarrow S1$ in all following rows of the matrix U .

In this way, the algorithm will work reliably even by multiple transformation.

(b) Skipping The Isolated Nodes

The algorithm marks by the coefficients $0^*, 0^{**}$ each node which becomes isolated in certain phase due to switching.

If the tested node has not zero coefficients in advance, its connection with the passive and active elements is tested. In case of finding first connection, we pass immediately to the test of next node. We can speed up this process considerably if the test is performed using the already compiled

coefficients. The algorithm can be described by following steps:

- we read node coefficients step by step from the first row of matrix $\mathbf{U}$, omitting the already read coefficients;
- we compare each coefficient with the coefficients of nodes to which the passive and active elements are connected.

In this way we test all connected nodes at the same time.

(c) Transformation Due To Nonregular Elements

Let us consider two types of ideal nonregular elements: the ideal differential OpAmp and the unity gain buffer.

Ideal differential OpAmp

Let the input nodes of the OpAmp be a_1 , a_2 and the output node be a_3 . Then the algorithm of the transformation of voltage coefficients a_1^* , a_2^* and a_3^* is the same as for the transformation due to switches.

The current coefficient of the output node has to be rewritten to zero on all positions of the matrices $\mathbf{U}_1$ and $\mathbf{U}_2$.

Buffer

Let the input/output node of the buffer be a_1/a_2 . Then the algorithm of the transformation of voltage coefficients a_1^* and a_2^* is the same as in the preceding case.

The current coefficient of the output node has to be rewritten to zero on all positions of the matrices $\mathbf{U}_1$ and $\mathbf{U}_2$.

(d) Rising Numbering Of The Coefficients

With the utilization of general properties of the voltage and current coefficients, following numbering algorithm has been developed and successfully verified:

1. We browse the matrix $\mathbf{U}_1/\mathbf{U}_2$ row by row. The voltage and current coefficients are classified separately. In each step, by now the highest coefficient S is stored. We set $S = 0$ before the first step.
2. If the coefficient in the row is greater than S , we increase the variable S by 1 and we set the value of this coefficient as S . In case this original

coefficient appears also in the other matrix rows, we replace also its value by S . If the coefficient in the actual row is less or equal to S , the values of S and of the coefficient remain unchanged. Then the algorithm passes to the next row.

5 CONCLUSION

Novel algorithm how to find the set of independent circuit variables in the switched circuits is presented. Computer simulator which utilizes this algorithm compiles only the minimum set of equations. The analysis is then fast and not memory consuming.

Proposed algorithm has been incorporated to the program SPIN for the analysis of general networks with the external switching [7].

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