

SECONDARY ROOT POLISHING TECHNIQUES IN COMPUTER CIRCUIT ANALYSIS

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Abstract: The so-called secondary root polishing techniques are described. They are used in the SNAP program to eliminate numerical inaccuracies during the zero/pole computation of the circuit functions of linear and linearized circuits.

Introduction

The analysis functions of some programs for the analysis of linear and linearized circuits are based on the circuit function zeros and poles. The SNAP program for the symbolic, semisymbolic, and numerical analyses belongs to this class [1]. The final analysis accuracy is then strongly dependent on the zero and pole finding algorithms. The question of accuracy of these algorithms is fundamental especially in some cases: partly in the analysis of large circuits of the order of about 100 and more, partly in circuits having a considerable spread of time constants and natural frequencies, or in circuits with multiple poles.

In principle, zeros/poles are computed by two different methods: either from the circuit matrix by the solution of generalized eigenvalue problem, or from the numerator/denominator polynomials by finding their roots. Generally, the first method is more accurate and - if possible – we always prefer it.

Nevertheless, in some cases of analysis, the zero/pole computation represents a numerically sophisticated problem, and even the very accurate algorithms can generate inadmissible errors. To decrease these errors, special techniques, called *root polishing*, are often used [2]. There are methods of iterative root improvement, which sometimes make it possible to decrease the error to a level which is given by the roundoff processor error. Let us call all these familiar methods *primary root polishing techniques*.

Let us assume that instead of a correct couple of poles $[-1, -1]$ we get after the primary polishing $[-1, -0.999999999999]$. This tiny error (10^{-10} %) causes an incorrect identification of two different poles instead of one twofold. Then the equation of the step response will be of quite a different structure than in the correct case.

The above example is a typical illustration of an error which can be eliminated by the so-called *secondary root polishing technique*.

Techniques of secondary polishing.

Parameters

$|Min|/|Max|, |Re/Im|, |Im/Re|, |Re|, |Im|$

control the zero and pole polishing using the method of zeroing the negligible components. This polishing removes the numerical inaccuracy, as given in following examples:

<u>set of roots before polishing</u>	<u>set of roots after polishing</u>
-3,2134586976e-34 +j2,3981289349e12	+j2,3981289349e12
-3,2134586976e-34 -j2,3981289349e12	-j2,3981289349e12
-8,8726534894e9+j1,99000298983	-8,8726534894e9
2,3399478723e-9	0

Description of the polishing process for the poles (for the zero, the procedure is equivalent):

1. Polishing using the zeroing of negligible components.

In the first step, the maximum absolute value MAX of all poles is computed.

Criterion $|Min|/|Max|$: Zeroing of poles with "too small" absolute values.

All the poles whose absolute values are smaller than

$$MAX * |Min|/|Max|$$

are set to zero.

From the physical point of view, the absolute values of roots represent the natural frequencies of first- or second-order subcircuits, which model the frequency response of analyzed circuit. That is why the $1/MAX$ gives the minimum of all natural frequencies, and $|Min|/|Max|$ gives the relationship between the minimum and maximum frequency. The default value

$$|Min|/|Max| = 10^{-5}$$

implies that all natural frequencies which exceed the minimum one by at least 5 decades are not considered in the frequency response.

In some cases, when the analyzed circuit has a larger spread of natural frequencies, $eps1$ could be decreased. For example,

$$|Min|/|Max| = 10^{-m}$$

defines the interval of m decades between the minimum and the maximum natural frequencies.

In the SNAP program, this parameter can be modified from 0 to 10m. The value 0 means that polishing is disabled.

Criteria $|Re/Im|$, $|Im/Re|$, $|Re|$, $|Im|$: Zeroing of negligible real or imaginary parts of poles.

Generally, every pole is a complex number $a+jb$.

<u>If</u>	<u>then</u>
$abs(a) < abs(b) * Re/Im $	$a=0$
$abs(b) < abs(a) * Im/Re $	$b=0$
$abs(a) < Re $	$a=0$
$abs(b) < Im $	$b=0$

From the physical point of view, the real parts of roots represent the reciprocal values of time constants τ of first- or second-order subcircuits which model the transient responses of analyzed circuit. The imaginary parts represent angular frequencies $\omega=2\pi/T$ of oscillations of the given subcircuits in their responses to quick changes in the exciting signals.

From this point of view, the inequalities given above represent the following relations:

$$\omega\tau = 2\pi \frac{\tau}{T} > \frac{1}{|Re/Im|}$$

$$\omega\tau = 2\pi \frac{\tau}{T} < |\text{Im}/\text{Re}|$$

$$\tau > \frac{1}{|\text{Re}|}$$

$$\omega < |\text{Im}|$$

Criterion $|\text{Re}/\text{Im}|$ ($|\text{Im}/\text{Re}|$) eliminates too large/small relationships τ/T . To set the implicit values 10 of both parameters, we replace all poles whose relationship τ/T is more than 15915 by their real parts, and all poles whose relationship τ/T is less than 15915 by their imaginary parts.

In some cases, e.g. in the analysis of resistive circuits, while modeling the influence of parasitic reactances, parameter $|\text{Re}/\text{Im}|$ should be decreased. On the contrary, analyzing losses of idealized reactance circuits may sometimes require decreasing of $|\text{Im}/\text{Re}|$.

Criterion $|\text{Re}|$ eliminates too large time constants τ , criterion $|\text{Im}|$ too small oscillation frequencies ω . Setting the implicit values 10m for both parameters, we replace all poles whose time constants τ exceed 100 seconds by their imaginary parts, and all poles whose natural frequencies are less than 0.01 rads/sec by their real parts.

For the circuits with the time constants larger than 100 seconds, parameter $|\text{Re}|$ has to be decreased.

For the circuits with extremely slow oscillations (<0,01 rads/sec), parameter $|\text{Im}|$ has to be decreased.

2. Identification of multiple roots.

Criterion **r12**: identification of multiple roots.

Identification of multiple roots is necessary for the partial fraction expansion of the s-domain circuit function, and for the subsequent inverse Laplace transform for finding equations of the transient and impulse responses

In the first step, all roots are sorted into groups with "almost the same" absolute values: Absolute values in each group are equal in the frame of relative error **r12**. Inside each group, the "coincide test" is performed: both the real and imaginary parts have to be the same in the frame of relative error **r12**. If the roots are identified as the identical, their real and imaginary parts are replaced by the mean values.

Parameter **r12 must be neither too large nor too small. For a large **r12**, two different roots can be identified as multiple. For a small **r12**, multiple roots can be identified as different. This can be dangerous especially for the complex roots, when the algorithm need not find their conjugate counterparts. In such cases, parameter **r12** has to be increased.**

Root polishing is finished by the algorithms which utilize parameters **m0** and **m9**.

3. Polishing chains of zeros and nines.

Criterion **m0**: identification of a string of **m0** zeros.

This root polishing method is explained in the following example:

<u>before polishing</u>	<u>after polishing</u>
-3.230000000047e-8	-3.23e-8

This algorithm tests the root's mantissa, whether a string of at least m_0 zeros is present on the right side of the decimal point. If it is found, all numbers to the right of this string are replaced by zeros.

The longer the string of zeros the greater the probability that the subsequent nonzero numbers are due to rounding errors.

The permitted values of m_0 are integers from 2 to 13.

Criterion m_9 : identification of a string of m_9 nines.

This root polishing method is explained in the following example:

<u>before polishing</u>	<u>after polishing</u>
-4.7999999999873e-4	-4.8e-4
0.9999999997389	1

This algorithm tests the root's mantissa, whether a string of at least m_9 nines is present on the right of the decimal point. If it is found, this string is replaced by zeros including all subsequent numbers, and a prospective order transfer is performed to the left-side.

The longer the string of nines the greater the probability that it is due to rounding errors.

The permitted values of m_9 are integers from 2 to 13.

Conclusion

The described methods of secondary root polishing have proved to be a useful complement of the well-known primary root polishing techniques. They are especially important for the identification of multiple roots.

These techniques are included in the SNAP network analyzer [1]. Controlling constants of the individual algorithms are user-accessible in the menu of global settings. The default setting ensures the proper root polishing for a wide class of analyzed circuits.

It is also possible to extend the secondary root polishing techniques beyond the area of root computation, and to combine them with the primary polishing.

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References

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